

EMBEDDINGS INTO THOMPSON'S GROUP V AND $co\mathcal{CF}$ GROUPS

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ABSTRACT. Lehnert and Schweitzer show in [21] that R. Thompson's group V is a co-context-free ($co\mathcal{CF}$) group, thus implying that all of its finitely generated subgroups are also $co\mathcal{CF}$ groups. Also, Lehnert shows in his thesis that V embeds inside the $co\mathcal{CF}$ group $QAut(\mathcal{T}_{2,c})$, which is a group of particular bijections on the vertices of an infinite binary 2-edge-colored tree, and he conjectures that $QAut(\mathcal{T}_{2,c})$ is a universal $co\mathcal{CF}$ group. We show that $QAut(\mathcal{T}_{2,c})$ embeds into V , and thus obtain a new form for Lehnert's conjecture. Following up on these ideas, we begin work to build a representation theory into R. Thompson's group V . In particular we classify precisely which Baumslag-Solitar groups embed into V .

1. INTRODUCTION

1.1. History and context. There has been a long historical interplay between classes of formal languages, and classes of groups. This connection was first made by Max Dehn, who in 1911 stressed the importance of several formal problems associated with group presentations, one of which was the word problem for groups (Given a group G , is there an algorithm which determines, in finite time, whether or not any given finite product of generators is trivial?). Note that when one looks at the set of "words" in the generators which are equivalent to the identity in the group one has specified a formal language. The more complicated this formal language is, the more complex an algorithm would have to be in order to positively answer Dehn's question.

Thus, as classes of languages become more complex, some "corresponding" classes of groups become wider, and thus harder to comprehend in a meaningful fashion. Ways to build these correspondences are through the word or co-word problems for groups, but other flavours of correspondence have also been seen (see, e.g., [2, 25, 19, 16, 17, 18]). To date, classifications of corresponding sets of groups (for classes of languages) only exist for very simple classes of languages, but the results on the group theory side are quite striking. In order to discuss this further, we need to give a definition.

Given a finitely generated group $G = \langle X \rangle$, one can define the *language of the word problem* to be the set of words

$$WP(G) = \{w \in F(X) \mid w \equiv_G 1\}.$$

Similarly, the *language of the co-word problem* is defined to be

$$coWP(G) = \{w \in F(X) \mid w \not\equiv_G 1\}.$$

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And now, let us state some results on the group theory side.

In 1972, Anisimov in [2] shows that for a finitely generated group G , $WP(G)$ is a regular language if and only if G is finite. The proof of this is not difficult, but the idea of building the correspondence in the first place represents quite a step forward. Later, in the early 1980's, a celebrated collection of papers of Muller and Schupp (relying on Dunwoody's accessibility theory - see [11, 24, 25]) show the following theorem.

Theorem 1 (Muller, Schupp). *Let G be a finitely generated group. Then, $WP(G)$ is a context-free language if and only if G is virtually free.*

A group G with $WP(G)$ a context free language is called a *context-free group* or a \mathcal{CF} group.

One thus sees that \mathcal{CF} groups provide a generalisation of finite groups from the point of view of computer science, as the context-free languages are one of the simplest generalisations of regular languages (the difference arises as the machines that are used to create context-free languages are directed labelled graphs which have a stack for memory, and can make transitions based on this changing stack, while the machines for regular languages are simply finite directed graphs with no form of memory other than their structure, see, e.g., [20] for an introduction to automata and formal language theory).

Muller and Schupp's results are in some sense the last complete classification of a class of groups corresponding to a class of languages. Thus, focus has shifted to the $co\mathcal{CF}$ groups.

The class $co\mathcal{CF}$ of *co-context-free groups* is defined to be finitely generated groups G such that $coWP(G)$ is context-free. This is a generalisation as the virtually free groups actually have *deterministic* context-free word problems, and thus their language of co-words is also a deterministic context-free language.

In [19] Holt, Rees, Röver and Thomas introduce the $co\mathcal{CF}$ groups. They show that the class $co\mathcal{CF}$ of all co-context-free groups is closed under taking:

- taking finite direct products,
- taking restricted standard wreath products with context-free top groups,
- passing to finitely generated subgroups
- passing to finite index overgroups.

In [19] there are also various conjectures about how other operations interact with the class $co\mathcal{CF}$. Also, there is a discussion about whether certain very specific groups can be $co\mathcal{CF}$ groups. Currently, it is conjectured that $\mathbb{Z}^2 * \mathbb{Z}$ and the Grigorchuk group Γ are not in $co\mathcal{CF}$, and that certain wreath products cannot be in $co\mathcal{CF}$.

Now, let $\mathcal{T}_{2,c}$ be the infinite binary 2-edge-colored binary tree (left edges = red, right edges = blue), and let $\text{QAut}(\mathcal{T}_{2,c})$ be the group of all bijections on the vertices of $\mathcal{T}_{2,c}$ which respect the edge and color relationships, except for at possibly finitely many locations. Lehnert in [23] shows the following results, amongst others.

Theorem 2 (Lehnert). *The group $\text{QAut}(\mathcal{T}_{2,c})$ is a $co\mathcal{CF}$ group, and there is an embedding from R. Thompsons group V into $\text{QAut}(\mathcal{T}_{2,c})$.*

We note that Lehnert and Schweitzer [21] prove that the Higman-Thompson groups $G_{r,s}$ are in $co\mathcal{CF}$, which also shows that $V = G_{2,1}$ is in $co\mathcal{CF}$.

In his dissertation, Lehnert also makes the following conjecture:

Conjecture 3 (Lehnert). The group $\text{QAut}(\mathcal{T}_{2,c})$ is a universal $co\mathcal{CF}$ group.

Thus, Lehnert conjectures that a group G is in $co\mathcal{CF}$ if and only if it is finitely generated and it embeds in $\text{QAut}(\mathcal{T}_{2,c})$.

It is the main focus of this paper to discuss Lehnert's conjecture.

1.2. Our results, and the ongoing discussion. Lehnert and Schweitzer in [22] ask the authors of [6] whether they thought one could embed $\text{QAut}(\mathcal{T}_{2,c})$ into R. Thompson's group V , as those authors had just shown that $\mathbb{Z}^2 * \mathbb{Z}$ fails to embed into V (supporting the conjecture of Holt, Röver, Rees and Thomas). This question eventually lead the current authors to our main result, below.

Theorem 4. *There is an embedding $\text{QAut}(\mathcal{T}_{2,c}) \hookrightarrow V$.*

That is, one can now re-state Lehnert's conjecture as:

Conjecture 5 (Lehnert). R. Thompson's group V is a universal $co\mathcal{CF}$ group.

Thus, if Lehnert's conjecture is true, a group will be in $co\mathcal{CF}$ if and only if it is finitely generated and it embeds as a subgroup of V .

Working to understand the class $co\mathcal{CF}$ better, we will now discuss some of what is known about the subgroups of R. Thompson's group V .

It is known that V contains many embedded copies of non-abelian free groups, and indeed, many free products of its subgroups. In the paper [6], the authors give some more specific results. They find conditions under which particular restricted wreath products and free products of subgroups of V actually can embed into V . Also, they show that $\mathbb{Z} * \mathbb{Z}^2$ does not embed into V , supporting the conjecture of Holt, Röver, Rees, and Thomas. In [10] Nathan Corwin adds to these results by showing that $\mathbb{Z} \wr \mathbb{Z}^2$ does not embed into V .

Three important questions that are often asked, pertaining to subgroup structure of a given group G , are a) to decide if non-abelian free groups embed into G , b) to decide whether surface groups embed into G , and c) to decide if the Baumslag-Solitar groups embed into G .

In [26] Röver shows that if n is a proper divisor of m , then the Baumslag-Solitar group $BS(m, n)$ does not embed into V . Röver's proof is based on using the fact observed by Higman (the corollary to Lemma 9.3 (see [15])) that non-torsion elements in the Higman-Thompson groups fail to have infinitely many roots. Using a different method, we extend Röver's result as follows.

In this paper we show that certain Baumslag-Solitar groups are $co\mathcal{CF}$ groups. In particular, we decide exactly which Baumslag-Solitar groups embed in R. Thompson's group V (and indeed, into the Higman-Thompson groups $G_{s,r}$).

Theorem 6. *Let $m, n \in \mathbb{Z} \setminus \{0\}$. Let $BS(m, n)$ be the corresponding Baumslag-Solitar group.*

- (1) *If $|m| \neq |n|$, then $BS(m, n)$ fails to embed in V .*
- (2) *If $|m| = |n|$, then there is an embedding of $BS(m, n)$ in V .*

In fact, the theorem above holds for all of the Higman-Thompson groups $G_{r,s}$ (not just $V = G_{2,1}$) using essentially the same proof we give.

At the time of this writing, Burillo, Cleary, Röver, and Stein are working on a survey of obstructions to finding embeddings into R. Thompson's group V , which will feature various arguments including an obstruction based on distortion of subgroups. One can use the distortion obstruction to give the above non-embedding

results for the Baumslag-Solitar groups, and in fact the argument is equivalent to our presented argument. We discuss this briefly in Section 5.

It is a question of Gromov as to whether the surface groups embed into all (word-)hyperbolic groups [4], and there has been much work on this question by the broader community. Amongst many results and partial results, one can highlight the general results of the Calegari school on stable commutator length, and in particular the work in [8, 9, 28] which shows amongst other things that in HNN extensions of free groups, amalgamating the base free group to an endomorphic embedded copy of itself, one can often find surface subgroups. The present authors have attempted to use these results to find surface groups (other than the torus and the Klein bottle groups) in V , but so far we have been unsuccessful. Thus, we ask the following question.

Question 7. Do the hyperbolic (closed) surface groups embed into R. Thompson's group V ?

We note in passing that it is known that the set of subgroups in V is closed in passing to finite index over-groups (see e.g. [6, 26]), thus all that is required is to find a copy of any single hyperbolic (closed) surface subgroup in V to see that they all embed.

Even if hyperbolic (closed) surface groups do not embed into R. Thompson's group V , it may still be the case that they are in the class $co\mathcal{CF}$, and so it is also natural to ask the following question.

Question 8. Are the hyperbolic (closed) surface groups in the class $co\mathcal{CF}$?

In response to an early draft of this article Daniel Farley in [13] describes a particular list of $co\mathcal{CF}$ groups arising from his constructions with Bruce Hughes, which might be used to disprove Lehnert's conjecture, if one can show that one of these groups fails to embed into V .

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2. DECOMPOSITION OF ELEMENTS OF $\text{QAut}(\mathcal{T}_{2,c})$

In this section we define two ways to describe an element of $\text{QAut}(\mathcal{T}_{2,c})$. The first form associates a minimal element of V with an element of $\text{QAut}(\mathcal{T}_{2,c})$, but is complicated by the intervention of two further bijections between finite subsets of $\{0, 1\}^*$, while the second way is simpler, and associates any element of $\text{QAut}(\mathcal{T}_{2,c})$ to a non-unique element of V and a bijection between two finite subsets of $\{0, 1\}^*$. We believe our first form is new while the second appears to be what is used by Lehnert

in his dissertation [23]. Still, the first form enables us to build an embedding of $\text{QAut}(\mathcal{T}_{2,c})$ into V in Section 4.

To set up the notation we will be using, we define by $\{0,1\}^*$ as the set of all finite words in the alphabet $\{0,1\}$ and by $\{0,1\}^\omega$ the set of infinite words in the alphabet $\{0,1\}$. The set $\{0,1\}^\omega$ corresponds to the boundary of the tree \mathcal{T}_2 and to the standard ternary Cantor set \mathfrak{C}_2 . If a word $a \in \{0,1\}^*$ is a *prefix* of a word w , we write $a \leq w$. If $a_1, a_2 \in \{0,1\}^*$ we write $a_1 \perp a_2$ if neither is a prefix of the other.

2.1. Building an element in V from an element in $\text{QAut}(\mathcal{T}_{2,c})$. Let $\tau \in \text{QAut}(\mathcal{T}_{2,c})$ be seen as a map $\tau : \{0,1\}^* \rightarrow \{0,1\}^*$. For any $w \in \{0,1\}^*$ we examine the pair $(w, w\tau)$ to find the largest common suffix s_w such that

$$\begin{aligned} w &= x_w s_w \\ w\tau &= y_w s_w \end{aligned}$$

for suitable $x_w, y_w \in \{0,1\}^*$ prefixes. We define $\Gamma_\tau := \{(x_w, y_w) \mid w \in \{0,1\}^*\}$.

Claim 9. Γ_τ is finite.

Proof. Since the map τ is a quasi-automorphism, there exists a level k in the domain tree below which the adjacency and color relations are respected by the action of τ . Assume that w_1 is any node below level k and let $w_2 := w_1 a$, for some $a \in \{0,1\}$. Since w_1 and w_2 are adjacent and below level k , we have $w_2 \tau = (w_1 a) \tau = (w_1 \tau) a$. This immediately extends to any word $\lambda \in \{0,1\}^*$, so that $(w_1 \lambda) \tau = (w_1 \tau) \lambda$. Since this argument holds for any w_1 below level k , this shows that there can be only finitely many elements in Γ_τ since every word below level k is a descendant of the finite set of words at level k . \square

Since the set Γ_τ is finite we can find a subset $M_\tau \subseteq \Gamma_\tau$ which is essential in the following sense: for every pair $(a, b) \in M_\tau$ there exist infinitely many words w such that $(x_w, y_w) = (a, b)$.

Recall that an *anti-chain* is a subset of a partially ordered set such that any two elements in the subset are incomparable and that an anti-chain is *complete* if it is maximal with respect to inclusion. By the definition of M_τ the sets of words

$$\mathcal{L}_{D_\tau} := \{a \in \{0,1\}^* \mid \exists b \in \{0,1\}^*, (a, b) \in M_\tau\}$$

and

$$\mathcal{L}_{R_\tau} := \{b \in \{0,1\}^* \mid \exists a \in \{0,1\}^*, (a, b) \in M_\tau\}$$

both form finite complete anti-chains for the poset $\{0,1\}^*$ ordered by prefix inclusion. The set \mathcal{L}_{D_τ} has the properties that if $a_1 \neq a_2 \in \mathcal{L}_{D_\tau} \subset \{0,1\}^*$ then $a_1 \perp a_2$ and for all sufficiently long words w in $\{0,1\}^*$, there is $a_w \in \mathcal{L}_{D_\tau}$ so that $a_w < w$, and similarly for \mathcal{L}_{R_τ} . Therefore, M_τ naturally determines a prefix code map on $\{0,1\}^\omega$ which is determined by two finite complete anti-chains of equal cardinality. Note that this is another way of defining element of R. Thompson's group V .

Furthermore, one easily sees that the construction of M_τ as above produces the bijection between the leaves of the unique minimal tree-pair representative for the particular element v_τ of V which M_τ determines. As any finite rooted subtree of \mathcal{T}_2 is determined by its leaves and vice-versa, we can identify the finite complete anti-chains $\mathcal{L}_{D_\tau}, \mathcal{L}_{R_\tau}$ with finite trees D_τ, R_τ and say that $(D_\tau, R_\tau, \sigma_\tau)$ is the minimal tree pair representative for the element v_τ that M_τ determines. We

write $v_\tau \sim (D_\tau, R_\tau, \sigma_\tau)$ in this case to emphasise that $(D_\tau, R_\tau, \sigma_\tau)$ is the minimal tree pair representing v_τ . (Note: we also write $v_\tau \sim (D, R, \sigma)$ if (D, R, σ) is any tree pair representing v_τ .)

Remark 10. There may exist a word $w \in \{0, 1\}^*$ such that there is a pair $(a, b) \in M_\tau$ with a prefix of w but $w \rightarrow w\tau$ is not determined by the pair (a, b) . This will lead us to consider a finite permutation also associated with τ , which we will discuss below in Subsection 2.3.

2.2. The finite bijection between internal nodes of the tree pair in v_τ . By a slight abuse of notation we identify the tree D_τ with the set of words corresponding to the nodes of τ strictly above the leaves. We make a similar identification for R_τ .

We now define a (non-canonical) bijection $b_\tau : D_\tau \rightarrow R_\tau$. For every word $w \in D_\tau \cap (R_\tau)\tau^{-1}$, we define $wb_\tau := w\tau$. Then we complete b_τ to a bijection by choosing and fixing a bijection between the sets $D_\tau \setminus (R_\tau)\tau^{-1}$ and $R_\tau \setminus (D_\tau)\tau$.

2.3. Finite permutation on a subset of the nodes of the tree R_τ . We observe that v_τ can be seen as a tree pair diagram or a map acting on $\{0, 1\}^* \setminus D_\tau$ as a prefix replacement map. Consider the following map

$$\tilde{v}_\tau = \begin{cases} b_\tau & \text{over } D_\tau \\ v_\tau & \text{over } \{0, 1\}^* \setminus D_\tau. \end{cases}$$

The map \tilde{v}_τ is an element of $\text{QAut}(\mathcal{T}_{2,c})$ and, by construction, it differs with τ on only finitely many vertices (which are possibly spread between D_τ and $\{0, 1\}^* \setminus D_\tau$). Thus the map $p_\tau := \tilde{v}_\tau^{-1}\tau$ is a permutation on finitely many vertices of $\mathcal{T}_{2,c}$.

2.4. Minimal decomposition of elements in $\text{QAut}(\mathcal{T}_{2,c})$. We are now able to write down the decomposition for $\tau \in \text{QAut}(\mathcal{T}_{2,c})$ that we were looking for. The following result is an immediate consequence of the discussion above.

Lemma 11 (Minimal decomposition). *For every $\tau \in \text{QAut}(\mathcal{T}_{2,c})$, there exists a permutation p_τ on finitely many vertices of $\mathcal{T}_{2,c}$ so that*

$$(2.1) \quad \tau = \tilde{v}_\tau p_\tau$$

where $\tilde{v}_\tau \in \text{QAut}(\mathcal{T}_{2,c})$ acts as an element of Thompson's group V beneath a suitable level and is a bijection on the finitely many nodes above such level.

We say that the decomposition of Lemma 11 is a *minimal decomposition* because the tree pair for the associated element of V is minimal.

Remark 12. As observed above, this decomposition is not unique and depends on how we choose to build the map b_τ . Nevertheless, there is always a way to create the decomposition in (2.1).

2.5. Disjoint decomposition form for elements in $\text{QAut}(\mathcal{T}_{2,c})$. It is possible to rewrite the form of Lemma 11 so that v_τ is represented by a tree pair (D_d, R_d, σ_d) whose domain tree D_d is the full subtree of \mathcal{T}_2 at depth k , for some k , while $b_\tau : X \rightarrow Y$ is a bijection from the set X of vertices of \mathcal{T}_2 of depth less than k , to the set Y of vertices of \mathcal{T}_2 which are above the leaves of R_d , and with p_τ equal to the identity map.

Lemma 13 (Disjoint decomposition). *For every $\tau \in \text{QAut}(\mathcal{T}_{2,c})$, there exists a map $d_\tau \in \text{QAut}(\mathcal{T}_{2,c})$ which acts as an element of Thompson's group V beneath a suitable level and is a bijection on the finitely many nodes above such level and such that*

$$(2.2) \quad \tau = d_\tau.$$

Proof. We use the same notation of the previous subsections and apply Lemma 11 to the element τ to rewrite it as $\tilde{v}_\tau p_\tau$. We then unreduce the tree pair diagram (D_τ, R_τ, σ) of v_τ to a new pair $(D'_\tau, R'_\tau, \sigma')$ where D'_τ is a full subtree chosen so that the leaves of each of the trees D'_τ and R'_τ are strictly below the set of vertices non-trivially acted upon by both b_τ and p_τ . We now define a map $b'_\tau : D'_\tau \rightarrow R'_\tau$ by $(t)b'_\tau := (t)\tilde{v}_\tau$ and the map $p'_\tau : R'_\tau \rightarrow R'_\tau$ by $(t)p'_\tau := (t)p_\tau$. By construction and Lemma 11, it is obvious that $(t)\tau = (t)b'_\tau p'_\tau$ for $t \in D'_\tau$.

We thus define $s_\tau : D'_\tau \rightarrow R'_\tau$ by $(t)s_\tau := b'_\tau p'_\tau$ and finally we define

$$d_\tau = \begin{cases} s_\tau & \text{over } D'_\tau \\ \tilde{v}_\tau & \text{over } \{0, 1\}^* \setminus D'_\tau. \end{cases}$$

It is immediate from our construction that $\tau = d_\tau$ and that d_τ is built as an element of V below some level and a bijection above such level. \square

We say that the decomposition of Lemma 13 is a *disjoint decomposition* because d_τ is described via a bijection s_τ and an element of Thompson's group V , that is \tilde{v}_τ restricted to the lowest vertices in $\{0, 1\}^* \setminus D'_\tau$ (which clearly define a tree). We sometimes refer to s_τ as the **bijection part** of d_τ and to element of the Thompson's group V given by restricting \tilde{v}_τ to the lowest vertices of $\{0, 1\}^* \setminus D'_\tau$ as the **V -part of d_τ** .

Remark 14. Given a map $\tau \in \text{QAut}(\mathcal{T}_{2,c})$ we can canonically define the **cutoff level of τ** as the smallest level k of the domain tree such that the map τ behaves as a colour and adjacency preserving automorphism on every vertex on every level $\geq k$. If $I(\tau)$ is the set of vertices in the domain of τ such that τ does not respect either the adjacency relation or the color relation, and ℓ is the largest level of any point inside $I(\tau)$, then $k = \ell + 1$.

We observe that the definition of the cutoff level depends only on τ , but one can also recover it via the support of p_τ . If we denote by $\text{Supp}(p_\tau)$ the support of p_τ and consider the set

$$Z(p_\tau) := (D_\tau \setminus (R_\tau)\tau^{-1}) \cup (R_\tau \setminus (D_\tau)\tau)\tau^{-1} \cup \text{Supp}(p_\tau)\tau^{-1}$$

and r is the largest level of any point in $Z(p_\tau)$, then $k = r + 1$. If k is the cutoff level of τ , we call the disjoint decomposition built via the full subtree of level k as the **cutoff disjoint decomposition of τ** . We observe that this decomposition is unique since the cutoff level is uniquely defined.

Lemma 15. *Any disjoint decomposition can be obtained by refining the cutoff disjoint decomposition.*

Proof. This is immediate from the proof of Lemma 13 and the definition of the cutoff level as any full subtree D'_τ used to build a disjoint decomposition must have depth greater or equal than k , the depth of the cutoff level. \square

3. AN EMBEDDING $V \hookrightarrow \text{QAut}(\mathcal{T}_{2,c})$

In this section we provide an embedding $\Theta : V \hookrightarrow \text{QAut}(\mathcal{T}_{2,c})$. Our embedding is similar in spirit to that Lehnert describes in his dissertation [23]. We provide this embedding as we were not able to directly verify the embedding Lehnert describes.

Given an element of V , we will embed it in $\text{QAut}(\mathcal{T}_{2,c})$, with support over the union of the set of words which begin with ‘0’ together with the set which contains only the empty word. No word beginning with ‘1’ will be moved by our embedding. Intuitively, our map will be what one gets if one associates V as acting on the ordered set $(0, 1/2]$ by interval exchange maps which exchange intervals of the form $(a, b]$ where a and b are dyadic rationals in the set $(0, 1/2]$ (where here, the root of $\mathcal{T}_{2,c}$ is corresponding to the value $1/2$, and our embedded copy of V is acting only on the left half of the interval $[0, 1]$ (fixing 0). Thus, our embedding really will be in the spirit of Lehnert’s embedding, we now formalise this discussion.

Rule for the injection: Given any finite rooted binary tree T , overlay the tree on $\mathcal{T}_{2,c}$ so that the root of T will be placed at the node 0 of $\mathcal{T}_{2,c}$. Refer to the embedded tree as T' . Let X' be the set of interior nodes of T' as a subset of the nodes of $\mathcal{T}_{2,c}$, and set $X := X' \cup \{\varepsilon\}$, that is, X' together with the empty node. Now associate a bijection ω_T from the leaves of T' to the nodes in X . We associate the leaves of T' to the nodes of X in left-to-right order (as seen in the tree where again, 0 means “left child” and 1 means “right child”). In particular, the rightmost leaf of T' is associated with the node ε .

Now, given an element $\alpha \in V$, let us describe the image $\alpha\Theta$ in $\text{QAut}(\mathcal{T}_{2,c})$. Suppose $\alpha \sim (D, R, \sigma)$. Then we can embed both D and R in $\mathcal{T}_{2,c}$ each as in the above paragraph to find embedded images D' and R' , remembering the associations from the leaves of these trees to the finite words in $\{0, 1\}^*$ which correspond to nodes of these trees (or the empty word) after they are embedded in $\mathcal{T}_{2,c}$. Now the permutation σ informs us how to move the maximal subtrees of the tree $\mathcal{T}_{2,c}$ rooted at the leaves of D' so they become trees rooted at the leaves of R' . If n is a node of $\mathcal{T}_{2,c}$ associated to a leaf of D' via ω_D^{-1} (so, n is either an interior node of D' or the root node of $\mathcal{T}_{2,c}$), then it should be mapped to the node $n\omega_D^{-1}\sigma\omega_R$ which is a node of $\mathcal{T}_{2,c}$ associated to a leaf of R' by the map ω_R .

We leave the following lemma as a technical exercise for the reader. As described below, this lemma enables one to verify that Θ is well defined (despite apparently depending on an initial choice of tree-pair for a given element of V). It then follows that Θ is a non-trivial group homomorphism.

Lemma 16. *Suppose that $\alpha \in V$ has $\alpha \sim (D_1, R_1, \sigma_1)$ and $\alpha \sim (D_2, R_2, \sigma_2)$, where D_2 is an elementary expansion of D_1 (that is, D_2 contains D_1 , and has exactly one extra caret). Then R_2 is an elementary expansion of R_1 as well, and the map Θ will send both representative tree pairs (D_1, R_1, σ_1) and (D_2, R_2, σ_2) to the same element in $\text{QAut}(\mathcal{T}_{2,c})$.*

The first consequence of Lemma 16 is that any two tree pairs representing the same element have the same image through Θ and so Θ is well-defined.

The second consequence is that the map Θ is actually a group homomorphism. One can see this as if $\alpha, \beta \in V$, then they induce specific permutations of the vertices in $\mathcal{T}_{2,c}$ which are independent of the tree pairs representing α, β (even though Θ is determined by these tree pairs). Thus, by expanding the tree pairs so that the range tree of α is identical to the domain tree of β , we see that the bijection of the

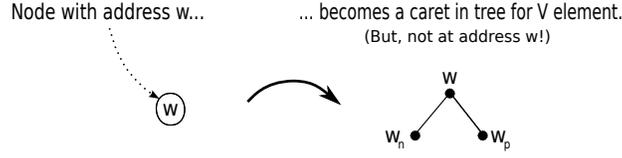


FIGURE 1. The replacement rule for nodes

vertices of $\mathcal{T}_{2,c}$ given by the composition of the maps $\alpha\Theta$ and $\beta\Theta$ is precisely the bijection given by the domain tree of α mapping to the range tree of β under the composition of the maps σ_α and σ_β (which is $(\alpha\beta)\Theta$).

We can observe that this homomorphism is non-trivial since applying it to the element of V that swaps the prefixes 0 with 1 already produces a non-trivial element of $\text{QAut}(\mathcal{T}_{2,c})$. Then simplicity of V implies that we have an embedding.

4. AN EMBEDDING $\text{QAut}(\mathcal{T}_{2,c}) \hookrightarrow V$

In this section we show the existence of a group homomorphism $\varphi : \text{QAut}(\mathcal{T}_{2,c}) \rightarrow V$ and, for $\tau \in \text{QAut}(\mathcal{T}_{2,c})$, we define it via the disjoint decomposition d_τ .

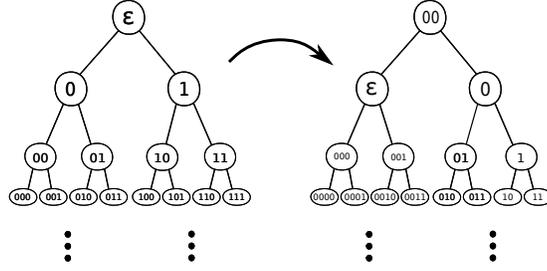
For every $\tau \in \text{QAut}(\mathcal{T}_{2,c})$, we apply Lemma 13 and obtain a disjoint decomposition of τ as an element $d_\tau \in \text{QAut}(\mathcal{T}_{2,c})$. We let D_{d_τ} and R_{d_τ} be the domain and range trees so that d_τ acts as an element of V on the leaves of D_{d_τ} (and below them) and as a bijection above such leaves.

We now apply the following construction to the trees D_{d_τ} and R_{d_τ} . We only explain it for the tree D_{d_τ} , the other being analogous. Assume that a vertex $w \in D_{d_\tau}$ has two children edges e_{left} and e_{right} and one parent edge e_{parent} in the tree. We replace w with a caret whose vertices are labeled by (w, w_n, w_p) , where w_n is the left child of w and w_p is the right child of w . We attach the former parent edge e_{parent} on top of w and we attach the former edges e_{left} and e_{right} below the vertex w_n . The vertex w_p has no children. We apply this construction to every vertex $w \in D_{d_\tau}$ with two exceptions:

- (1) the root vertex ε , for which there is no parent edge and so we only attach the two edges below the left child ε_n of ε .
- (2) Any leaf vertex w , to which we attach the two children edges with terminal vertices w_n and w_p .

The result of this construction is shown in figure 4. We denote the two trees we have just constructed by \widehat{D}_{d_τ} and \widehat{R}_{d_τ} . We say that all leaves of the form w_n are the **n -leaves** and the leaves of the form w_p are the **p -leaves**.

If $\mathcal{L}(T)$ defines the set of leaves of a finite binary tree T , we observe that the map d_τ induces a bijection $\sigma_{d_\tau} : \mathcal{L}(\widehat{D}_{d_\tau}) \rightarrow \mathcal{L}(\widehat{R}_{d_\tau})$ which we now construct. By definition, the map σ_{d_τ} acts on a leaf $w_n \in \mathcal{L}(\widehat{D}_{d_\tau})$ by sending it to the leaf $d_\tau(w)_n \in \mathcal{L}(\widehat{R}_{d_\tau})$ and acts on a leaf $w_p \in \mathcal{L}(\widehat{D}_{d_\tau})$ by sending it to a suitable leaf $t_p \in \mathcal{L}(\widehat{R}_{d_\tau})$ if the map d_τ sends its parent vertex w (seen as a vertex in D_{d_τ}) to a vertex t (seen as a vertex in R_{d_τ}). This defines a permutation σ_{d_τ} on the n -leaves (which comes from the permutation of the V -part of d_τ) and on the p -leaves (which comes from the bijection part of d_τ).



becomes...

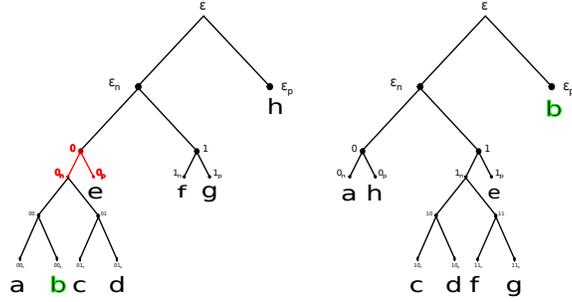


FIGURE 2. The replacement rule for nodes

We define $\varphi(\tau) := (\widehat{D}_{d_\tau}, \widehat{R}_{d_\tau}, \sigma_{d_\tau}) \in V$ (see figure 4). We need to verify that such map is well defined as it relies on a choice of a disjoint decomposition.

Lemma 17. *The map $\varphi : \text{QAut}(\mathcal{T}_{2,c}) \rightarrow V$ defined above is well-defined.*

Proof. By Lemma 15, any disjoint decomposition d_τ is built by starting from the cutoff disjoint decomposition and expanding the domain full subtree of the V -part coincides with D_{d_τ} . Any disjoint decomposition represents exactly the same map as τ , but the way it is written out is slightly different, even if the final outcome is the same function. Any two disjoint decomposition can be obtained from the cutoff one using unreductions of the V -part, thus one can go from one disjoint decomposition to another via a sequence of reductions and unreductions of the V -part. Hence, it is sufficient to prove our claim in the case of two disjoint decompositions f_τ and g_τ so that the tree pair defining the V -part of g_τ is obtained by an unreduction of a single level, that is adding a caret to each leaf in the domain and the range tree of the V -part of f_τ .

Let w be one of the leaves of f_τ to which we added two new children $w0$ and $w1$. By construction, the domain (respectively, range) tree of $\varphi(g_\tau)$ is obtained by adding the construction of figure 4 to the leaf w_n (respectively, adding the same construction to $f_\tau(w)_n = g_\tau(w)_n$). Moreover, the vertices $w_p, w0_p, w1_p$ are mapped

in an order preserving way to the vertices $f_\tau(w)_p = g_\tau(w)_p, g_\tau(w0)_p, g_\tau(w1)_p$. Observe that, as we added the exact same construction of figure 4 to the vertices w_n and $f_\tau(w)_n$ while building the tree diagram of $\varphi(g_\tau)$, we can immediately reduce it. We can do this for each and every of the carets we added to the V -part of f_τ . Hence, it is immediate that $\varphi(g_\tau) = \varphi(f_\tau)$. \square

Theorem 18. *The map $\varphi : \text{QAut}(\mathcal{T}_{2,c}) \rightarrow V$ is an injective homomorphism.*

Proof. *The map φ is a group homomorphism.* Let $\tau, \lambda \in \text{QAut}(\mathcal{T}_{2,c})$ and let f_τ and g_λ be disjoint decompositions built so that the range tree R_{f_τ} contains the full subtree constituting the domain tree D_{g_λ} . We can now unreduce the domain tree D_{g_λ} to make it become equal to R_{f_τ} and we unreduce R_{g_λ} accordingly to get a new tree R . We call $\bar{\sigma}_{g_\lambda}$ the bijection on the vertices which one obtains from σ_{g_λ} after this unredution. We define $g'_\lambda \in \text{QAut}(\mathcal{T}_{2,c})$ as the map which has a V -part defined by $(R_{f_\tau}, R, \bar{\sigma}_{g_\lambda})$ and a bijection part which is given by g_λ on all vertices above the set $\mathcal{L}(R_{f_\tau})$. Therefore, $g'_\lambda = g_\lambda$ as maps and, by a slight abuse of notation, we still say that g'_λ is a “disjoint” decomposition of λ as in this proof we only need the requirement that g'_λ is an automorphism on every vertex below $\mathcal{L}(D_{g'_\lambda})$. Therefore we have

$$(D_{f_\tau}, R_{f_\tau}, \sigma_{f_\tau})(D_{g_\lambda}, R_{g_\lambda}, \sigma_{g_\lambda}) = (D_{f_\tau}, R_{f_\tau}, \sigma_{f_\tau})(R_{f_\tau}, R, \bar{\sigma}_{g_\lambda}) = (D_{f_\tau}, R, \sigma_{f_\tau} \bar{\sigma}_{g_\lambda})$$

By construction, the tree pair $(D_{f_\tau}, R, \sigma_{f_\tau} \bar{\sigma}_{g_\lambda})$ has a domain tree which is a full subtree and is deep enough so that the action on all levels below is given by the composition of the V -parts of f_τ and of g_λ . This implies that the tree pair $(D_{f_\tau}, R, \sigma_{f_\tau} \bar{\sigma}_{g_\lambda})$ constitutes the V -part of a disjoint decomposition for the element $\tau\lambda$ (which we denote by $d_{\tau\lambda}$), where the associated bijection part can be computed by composing the bijections occurring within the trees $D_{f_\tau}, R_{f_\tau}, R$.

Now observe that $\varphi(d_{\tau\lambda})$ is obtained by expanding the trees D_{f_τ} and R and applying the $d_{\tau\lambda}$ to the n -leaves and the p -leaves. We obtain the tree pair $(\widehat{D}_{f_\tau}, \widehat{R}, \sigma) = \varphi(d_{\tau\lambda})$.

We consider now the tree pairs $\varphi(f_\tau)$ and $\varphi(g'_\lambda) = \varphi(g_\lambda)$. Since the range tree of the V -part of g'_λ is R , it is clear that the range tree of $\varphi(g'_\lambda)$ is equal to \widehat{R} . Therefore, the domain tree of $\varphi(d_{\tau\lambda})$ and of $\varphi(f_\tau)\varphi(g'_\lambda)$ is equal and the corresponding range trees coincide.

We observe that the bijection on the n -leaves of $\varphi(d_{\tau\lambda})$ and $\varphi(f_\tau)\varphi(g'_\lambda)$ is exactly the same because the V -part of $d_{\tau\lambda}$ was built by composing the tree pairs of the V -parts of f_τ and g'_λ . As we are dealing with disjoint decompositions, the actions above $\mathcal{L}(D_{f_\tau})$ and $\mathcal{L}(D_{g'_\lambda})$ do not affect the V -parts and so the bijections on the p -leaves of $\varphi(f_\tau)$ and of $\varphi(g'_\lambda)$ are exactly the bijections appearing in the bijection parts of f_τ and g'_λ . Thus the bijection on the p -leaves of $\varphi(f_\tau)\varphi(g'_\lambda)$ is determined by the composition of the bijection parts in f_τ and g'_λ which is equal to the bijection part of $d_{\tau\lambda}$. Therefore the bijection on the p -leaves of $\varphi(f_\tau)\varphi(g'_\lambda)$ is equal to that of $\varphi(d_{\tau\lambda})$. By putting everything together and using Lemma 17 we deduce that $\varphi(d_{\tau\lambda}) = \varphi(f_\tau)\varphi(g'_\lambda) = \varphi(f_\tau)\varphi(g_\lambda)$.

The map φ is injective. Let d_τ be the cutoff decomposition form for $\tau \in \ker \varphi$. The domain and the range tree of $\varphi(d_\tau)$ must be the equal and the permutation on every leaf is the identity permutation. Therefore the V -part of d_τ is the identity element and the bijection part is the identity map, that is d_τ is the identity map on $\mathcal{T}_{2,c}$ and so $\ker \varphi = \{1_{\text{QAut}(\mathcal{T}_{2,c})}\}$. \square

One observes that by composing the maps Θ and φ we can produce large families of embeddings of the type $V \rightarrow V$ and $\text{QAut}(\mathcal{T}_{2,c}) \rightarrow \text{QAut}(\mathcal{T}_{2,c})$. It is commonly known that V contains many copies of itself, however it is a question whether or not the embeddings produced here by multiple compositions of Θ and φ are of independent interest.

5. EMBEDDING BAUMSLAG-SOLITAR GROUPS IN V

In this section, we prove that the Baumslag-Solitar groups

$$BS(m, n) = \langle a, b \mid b^{-1}a^mb = a^n \rangle$$

fail to embed in R. Thompson's group V whenever $|m| \neq |n|$ and that they do embed otherwise. Röver observes that the following result is a consequence of Higman's work in [15].

Theorem 19 (Röver, [26]). *If n is a proper divisor of m , then the group $BS(m, n)$ does not embed in Thompson's group V .*

Recall that Farb and Franks have embedding results in groups related to Thompson groups. More precisely, they show the following

Theorem 20 (Farb-Franks, [12]). *Let m, n be positive integers.*

- (1) *If $m > n \geq 1$ the group $BS(m, n)$ embeds in the group of orientation-preserving analytic diffeomorphisms $\text{Diff}_+^\omega(\mathbb{R})$ and also inside the groups of orientation-preserving homeomorphisms $\text{Homeo}_+(S^1)$ and $\text{Homeo}_+(I)$.*
- (2) *If $m > n > 1$, the group $BS(m, n)$ does not embed into $\text{Diff}_+^2(I)$. If n does not divide m , the group $BS(m, n)$ does not embed into $\text{Homeo}_+(S^1)$.*

Taking the Farb and Franks result together with the Röver result, we have some evidence that most Baumslag-Solitar groups do not embed in T , and probably also in V . Thus, it is somewhat natural that only if $|m| = |n|$ should one expect that $BS(m, n)$ might embed in V (well, in R. Thompson's group T at the least). In this section, we show that these indications do not mislead.

The key idea behind the main result of this section is the following:

Lemma 21. *Let $v \in V$ be non-torsion, and r and s be integers. Then whenever $(v^r)^w = v^s$ for some $w \in V$, we have that $|r| = |s|$.*

The essence of the argument below will be clear to any reader who has digested the material on revealing pairs for elements of V (for instance, as in [5], which has an expository section written to explain these objects). Note that Higman in [15] introduces the technology of normal and semi-normal forms and later in [7] Brin introduces the analogous technology of revealing pairs, which we have chosen to use here.

Proof. If two elements $x, y \in V$, we use the conjugation notation $x^y = y^{-1}xy$. Assume that $v, w \in V$, v is not torsion, and that there are integers r and s so that $(v^r)^w = v^s$. We will now show that $|r| = |s|$. Note that if $r = 0$ the lemma result is immediately true, as $v^0 = 1_V$. Hence, we will assume below that neither r nor s is zero.

As v is not torsion, by an extension of an argument of Brin in [7] there is a minimal positive integer m so that, if we set $\alpha := v^m$, then $\langle \alpha \rangle$ is an infinite cyclic group embedded in V with the property that this group acts with no non-trivial

finite orbits on the Cantor set \mathfrak{C}_2 . For example, a possibility for m is the least common multiple of the set of lengths of all finite periodic orbits, as there are only finitely many such lengths. Now, the element α admits a finite set of points $I(\alpha)$, which we will call the *important points of α* (following [5]), consisting of the repelling and attracting points in the Cantor set under the action of $\langle \alpha \rangle$. For each point in the set $I(\alpha)$, it is the case that α restricted to some small interval U_p containing p is an affine map which fixes exactly the point p , where the slope of this map is 2^{s_p} , for some s_p a fixed non-zero integer. Now, as $(v^r)^w = v^s$ we see that

$$(\alpha^r)^w = ((v^m)^r)^w = ((v^r)^w)^m = (v^s)^m = (v^m)^s = \alpha^s$$

as well.

For any integer $u \in \mathbb{Z} \setminus \{0\}$, consider the finite set of logarithms of derivatives

$$\mathcal{S}_u = \left\{ \log_2 \frac{d}{dt} \alpha^u \Big|_{t=x} \mid x \in I(\alpha) \right\}.$$

Observe that $\mathcal{S}_u = u \cdot \mathcal{S}_1 = \{uv \mid v \in \mathcal{S}_1\}$. The equation above implies that w sends \mathcal{S}_r to \mathcal{S}_s . It is straightforward to verify that the finite set of slopes of $(\alpha^r)^w$ on $I(\alpha)$ is exactly equal to \mathcal{S}_r . Therefore

$$r \cdot \mathcal{S}_1 = \mathcal{S}_r = \mathcal{S}_s = s \cdot \mathcal{S}_1$$

and thus, if $k = \max\{|v| \mid v \in \mathcal{S}_1\} > 0$, one has

$$|r|k = \max\{|rv| \mid v \in \mathcal{S}_1\} = \max\{|sv| \mid v \in \mathcal{S}_1\} = |s|k$$

By the cancellation law, we have $|r| = |s|$. □

Remark 22. We would like to make a small historical comment. The core idea behind the proof of Lemma 21 is that the product of the set of slopes of the affine restrictions of v in a small neighbourhood of an orbit of a repelling or attracting periodic point (of \mathfrak{C}_2 associated with the action of $\langle v \rangle$) is an invariant of conjugacy in V of v . (See, e.g., [3, 14, 15, 27] for solutions of the conjugacy problem for V). One can think of the result as a total “speed” along an orbit, and it is useful in many ways. For instance, this product-of-slopes calculation for an orbit occurs in section 7 of [5] to analyse element centralisers. As described later in this section, it also appears in [5] to show that all cyclic subgroups are undistorted in V .

We are now ready to give a proof of Theorem 6.

Theorem 6. *Let $m, n \in \mathbb{Z} \setminus \{0\}$. Let $BS(m, n)$ be the corresponding Baumslag-Solitar group.*

- (1) *If $|m| \neq |n|$, then $BS(m, n)$ fails to embed in V .*
- (2) *If $|m| = |n|$, then there is an embedding of $BS(m, n)$ in V .*

Proof. Part (1) follows immediately by Lemma 21, since the element $a \in BS(m, n)$ is non-torsion.

For part (2) we consider the following construction. Let $e \in \{-1, 1\}$ and consider the group $BS(m, em)$. We now consider the diagonal embedding of $BS(m, em)$ into the direct product $(\mathbb{Z} * (\mathbb{Z}/m\mathbb{Z})) \times (\mathbb{Z} \rtimes_{\varphi} \mathbb{Z})$. The embedding of $BS(m, em)$ in the two coordinates is seen as follows:

- the left homomorphism is given by modding out by the common subgroup $\langle a^m \rangle$ inside $BS(m, em)$,

- the right homomorphism is the homomorphism from $BS(m, em) = \langle a, b \rangle$ to $BS(1, e) = \langle a', b' \rangle \cong (\mathbb{Z} \rtimes_{\varphi} \mathbb{Z})$, where $a \mapsto a'$ and $b \mapsto b'$.

It is straightforward to see that the intersection of the kernels of the two coordinates is trivial. The authors learned of the nice construction detailed above from Yves de Cornulier [1].

Now, the groups \mathbb{Z} and $\mathbb{Z}/m\mathbb{Z}$ are demonstrable groups for V , so Theorem 1.4 in [6] implies that $\mathbb{Z} * (\mathbb{Z}/m\mathbb{Z})$ is a subgroup of V . (A group is *demonstrable for V* if it admits a demonstrative embedding in V . A subgroup of V is *demonstrative* if there is a basic open set in the Cantor set \mathfrak{C}_2 which is moved entirely off of itself by any non-trivial element of the acting group. See Bleak and Salazar-Díaz [6] for more details.)

Separately, we need to discuss the following two cases:

- the group $BS(1, 1)$, which is isomorphic to \mathbb{Z}^2 is a well-known subgroup of V , and
- the group $BS(1, -1)$ which is the Klein bottle group.

Noting that there is a double cover of the torus to the Klein bottle, the Klein bottle group is a finite extension of the group $\mathbb{Z}^2 \leq V$ and therefore $BS(1, -1)$ is also a subgroup of V (by Corollary 1.3 of Bleak and Salazar-Díaz [6], the set of isomorphism classes of subgroups of V is closed under extensions by finite groups).

Therefore, $BS(m, em)$ embeds into the group $(\mathbb{Z} * (\mathbb{Z}/m\mathbb{Z})) \times (\mathbb{Z} \rtimes_{\varphi} \mathbb{Z})$ which is a subgroup of $V \times V$ which, in turn, is a subgroup of V . \square

We are grateful to José Burillo who pointed out that one can sometimes use an argument based on distortion to show certain non-embedding results. Indeed, this type of argument can be used in our context to show the main non-embedding results of Theorem 6. Here is how the argument runs.

First, recall Theorem 1.3 in [5] which says that all cyclic groups are undistorted in V . Now recall that when $|m| \neq |n|$, the group $BS(m, n)$ has distorted cyclic subgroups. Now we observe that if $BS(m, n)$ were a subgroup of V and $v \in BS(m, n)$ were an element of infinite order, then the distortion of $\langle v \rangle$ in V would have to be at least as much as it is in $BS(m, n)$ and that yields a contradiction, therefore implying that for such m and n , $BS(m, n)$ does not embed in V .

We observe that this argument is fundamentally equivalent to the one given in the proof of Lemma 21; to prove cyclic groups are undistorted in V , one measures the ‘speed’ of the elements in the cyclic subgroup, near to their attracting and repelling orbits.

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